

# MacWilliams Theorem for Linear Codes with Group Actions

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## Abstract

We prove MacWilliams theorem for linear codes with finite group actions. When acting group is trivial, our result becomes the ordinary MacWilliams theorem.

**Key words:** group, linear code, dual code, weight enumerator, MacWilliams identity.

## 1 Introduction and Summary

Yoshida [3] has given a version of the MacWilliams theorem [2] for codes with group action. In this paper we establish another version of the MacWilliams theorem. Our result seems to be a special case of Yoshida's. But we can not prove this.

Let  $V$  be the vector space  $\mathbf{F}_q^n$ , where  $\mathbf{F}_q$  is the field with  $q$  elements. From now on we assume that  $G$  is a finite permutation group on the coordinates of  $V$  and  $|G|$  is prime to  $q$ . Then we can define a natural action of  $G$  on  $V$  as follows: If  $\mathbf{v} = (v_1, \dots, v_n)$  and  $g \in G$ , we let  $\mathbf{v}g = (x_1, \dots, x_n)$ , where for  $i = 1, \dots, n$ ,  $x_i = v_{ig^{-1}}$ . In this way  $V$  becomes an  $FG$ -module. A  $G$ -code is an  $FG$ -submodule of  $V$ . As in [1], the operator  $\theta$  is defined by

$$\theta = \frac{1}{|G|} \sum_{g \in G} g.$$

Here we note that  $C_V(G) = V\theta$  and  $\theta^t = \theta$  (see [1]).

Let  $C_1, \dots, C_t$  be the orbits of the coordinates of  $V$  under the action of  $G$ . Let  $m_i$  be the orbit length of  $C_i$ . Define  $\overline{C}_i$  as the vector of  $V$  which has 1 as its entry for every

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point of  $C_i$  and 0 elsewhere. (This definition of the  $\overline{C}_i$ 's is slightly different from that in the proof of Theorem 4.3 in [1]). Then each of  $\overline{C}_1, \dots, \overline{C}_t$  is in  $V\theta$  and every element  $\mathbf{u}$  of  $V\theta$  is of the form

$$\mathbf{u} = \sum_{i=1}^t x_i \overline{C}_i.$$

This basis  $\{\overline{C}_1, \dots, \overline{C}_t\}$  of  $V\theta$  is a key to our proof. The  $G$ -weight of a vector  $\mathbf{u} = \sum_{i=1}^t x_i \overline{C}_i \in V\theta$  denoted  $wg(\mathbf{u})$  is defined as the number of non-zero  $x_i$ . So if  $G$  consists of the identity element,  $e$ , alone, then the  $G$ -weight  $wg(\mathbf{u})$  of a vector  $\mathbf{u}$  is the ordinary weight  $|\mathbf{u}|$ . For vectors  $\mathbf{a} = \sum_{i=1}^t a_i \overline{C}_i$ ,  $\mathbf{b} = \sum_{i=1}^t b_i \overline{C}_i$  of  $V\theta$ , an inner product  $(\mathbf{a}, \mathbf{b})_G$  of  $\mathbf{a}$  and  $\mathbf{b}$  is defined by

$$(1) \quad (\mathbf{a}, \mathbf{b})_G = a_1 b_1 + \dots + a_t b_t.$$

Let  $D$  be a vector subspace of  $V\theta$ .  $D_G^\perp$  is the dual of  $D$  in  $V\theta$  with respect to the inner product (1). (Notice that if  $G$  consists of the identity element,  $e$ , alone, then  $D_{\{e\}}^\perp$  is the ordinary dual  $D^\perp$  of  $D$  in  $V$ .)

We describe a weight enumerator of a vector subspace  $D$  of  $V\theta$ . The weight enumerator  $W_D(x, y)$  of  $D$  is defined by

$$W_D(x, y) = \sum_{\mathbf{u} \in D} x^{t-wg(\mathbf{u})} y^{wg(\mathbf{u})}.$$

Clearly if  $G$  is trivial, that is,  $G = \{e\}$ , then this weight enumerator becomes the ordinary weight enumerator. We shall prove the following:

**Theorem 1** *If  $C$  is a  $G$ -code, then*

$$W_{C^\perp\theta}(x, y) = \frac{1}{|C\theta|} W_{C\theta}(x + (q-1)y, x - y).$$

If  $G$  is trivial, that is,  $G = \{e\}$ , then our Theorem is the ordinary MacWilliams theorem [2, p. 146].

For notation and terminology, we shall refer the following book and paper: [2] for coding theory; [3] for codes with group action.

## 2 Proof of Theorem

In order to prove Theorem we need the following proposition.

**Proposition 1** *Let  $V$  be the vector space  $\mathbf{F}_q^n$ . Assume that  $G$  is a finite permutation group on the coordinates of  $V$  and  $|G|$  is prime to  $q$ . If  $C$  is a  $G$ -code and*

$$\theta = \frac{1}{|G|} \sum_{g \in G} g,$$

then

$$(C\theta)^\perp = \ker \theta + C^\perp\theta.$$

**Proof** See the proofs of Theorem 4.2 and Corollary 1 in [1].  $\square$

We shall prove Theorem. If  $\mathbf{x} = \sum_i x_i \bar{C}_i \in C\theta$  and  $\mathbf{y} = \sum_i y_i \bar{C}_i \in C^\perp\theta$ , by Proposition 1 we have

$$0 = (\mathbf{x}, \mathbf{y}) = \sum_i m_i x_i y_i = (\mathbf{x}, \mathbf{y}')_G,$$

where  $\mathbf{y}' = \sum_i m_i y_i \bar{C}_i$ . From this it follows that

$$(2) \quad (C\theta)_G^\perp \supseteq (C^\perp\theta)M,$$

where

$$M = \text{diag}(\underbrace{m_1, \dots, m_1}_{m_1 \text{ times}}, \underbrace{m_2, \dots, m_2}_{m_2 \text{ times}}, \dots, \underbrace{m_t, \dots, m_t}_{m_t \text{ times}}).$$

We shall show that

$$(3) \quad (C\theta)_G^\perp = (C^\perp\theta)M.$$

By Proposition 1, we have

$$(4) \quad \dim C^\perp\theta = \dim (C\theta)^\perp - \dim \ker \theta.$$

From linear algebra theory,

$$(5) \quad \dim V = \dim V\theta + \dim \ker \theta,$$

$$(6) \quad \dim V = \dim (C\theta)^\perp + \dim C\theta,$$

$$(7) \quad \dim V\theta = \dim (C\theta)_G^\perp + \dim C\theta.$$

From (4), (5), (6) and (7) we see that

$$(8) \quad \dim (C\theta)_G^\perp = \dim (C^\perp\theta).$$

Since  $M$  is a non-singular matrix, we have

$$(9) \quad \dim C^\perp\theta = \dim (C^\perp\theta)M.$$

From (2), (8) and (9) it follows that

$$(C\theta)_G^\perp = (C^\perp\theta)M.$$

Here notice that MacWilliams theorem [2, p. 146] for the ordinary weight enumerator of the code  $C\theta$  in  $V\theta$  holds in this case, too.

Now we shall finish the proof of Theorem. By MacWilliams theorem and (3), we obtain the following:

$$(10) \quad W_{(C^\perp\theta)_M}(x, y) = \frac{1}{|C\theta|} W_{C\theta}(x + (q-1)y, x-y).$$

Since  $W_{(C^\perp\theta)_M}(x, y) = W_{C^\perp\theta}(x, y)$ , it follows from (10) that

$$W_{C^\perp\theta}(x, y) = \frac{1}{|C\theta|} W_{C\theta}(x + (q-1)y, x-y).$$

□

**Remark.** Generalizing a result of Thompson, Hayden [1] has proved the following proposition.

**Proposition 2** *Using the notation of Proposition 1, then with an appropriate orthonormal base for  $V\theta$ , (extending  $\mathbf{F}_q$  if necessary) we have where  $(C\theta)_{V\theta}^\perp$  is the dual in terms of this basis*

$$(C\theta)_{V\theta}^\perp = C^\perp\theta.$$

So our result (3) is a generalization of Proposition 2 in a sense.

### References

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